

Relationship between communication and quantum metrology

arxiv:1603.00472

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03.03.2016
RAQM 2016

Communication Theory

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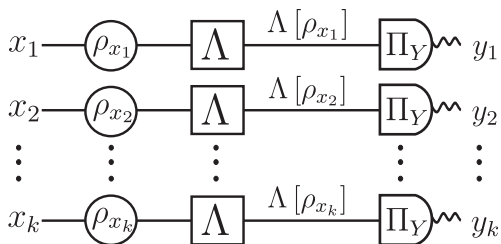
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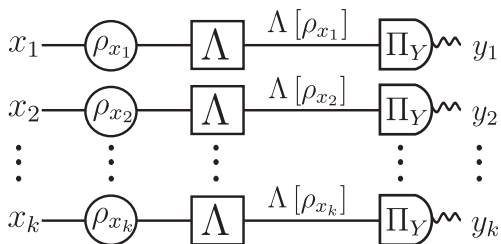
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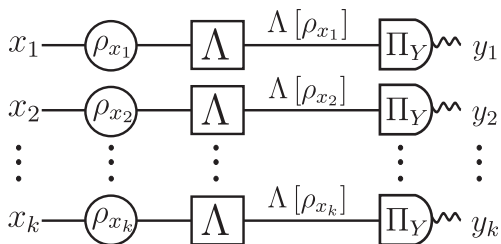
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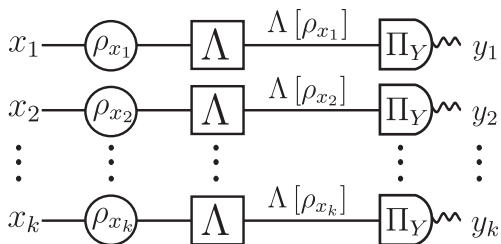
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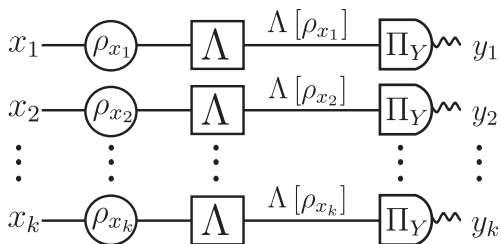
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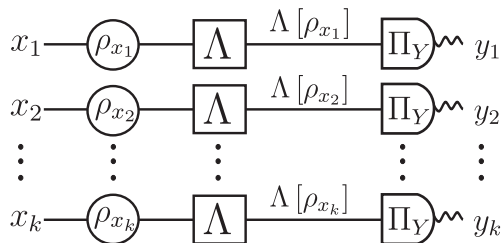


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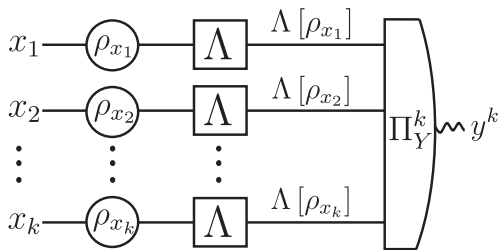
Superadditivity

- General scheme



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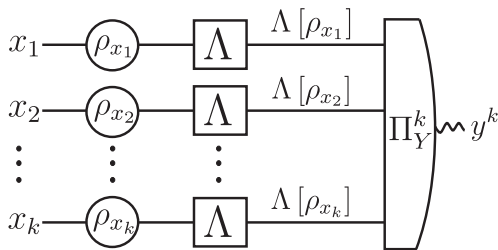
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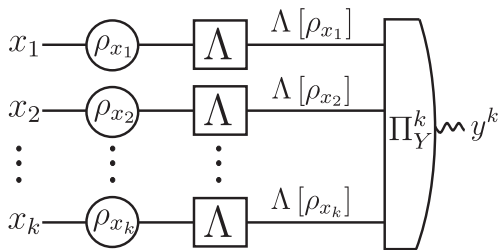
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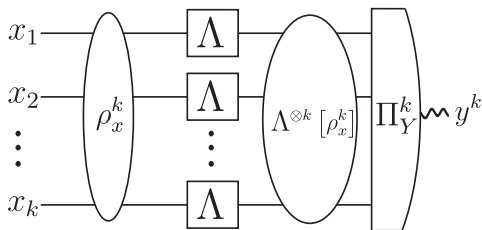
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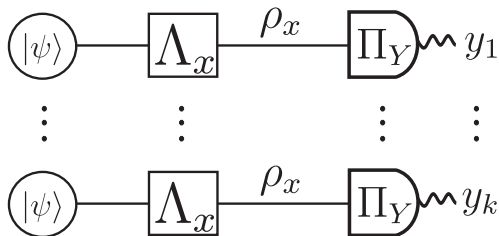
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- Input superadditivity: use entangled states \rightarrow
 $C^{(k,\infty)} \geq C^{(1,\infty)}$ (Hastings, Nat. Phys. (2009)).

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- We want to measure some parameter x with the best possible precision $\Delta x^2 = \int p(y|x)(\tilde{x}(y) - x)^2 dy$.

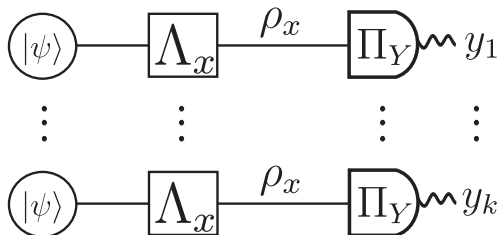
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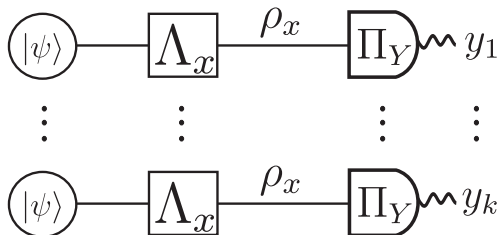
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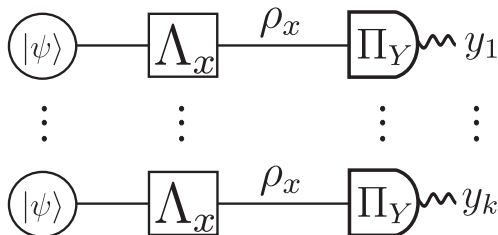
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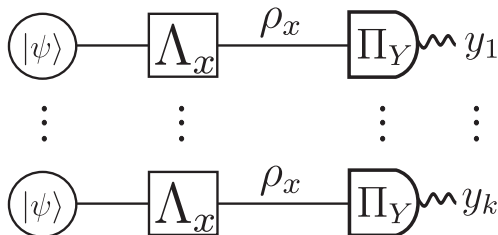


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- Bound on precision optimized over all unbiased estimators.

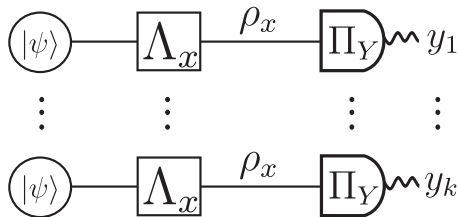
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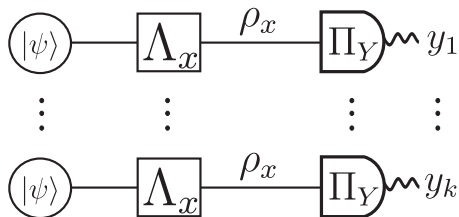
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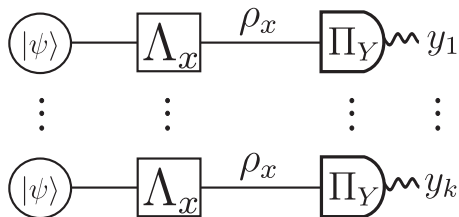


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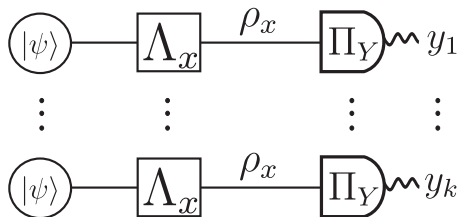
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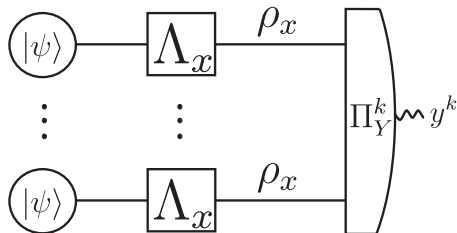
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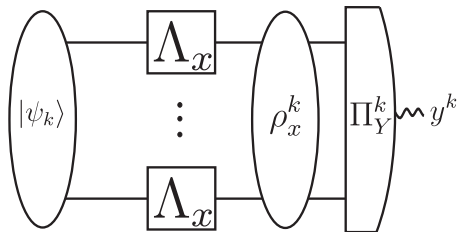
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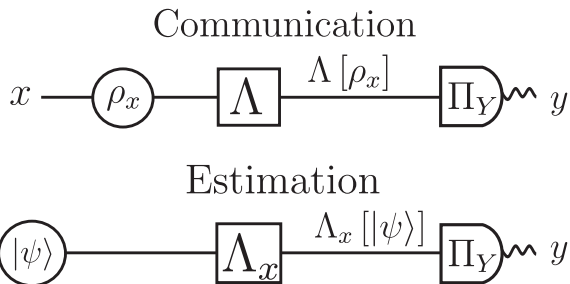
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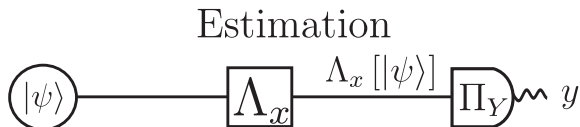
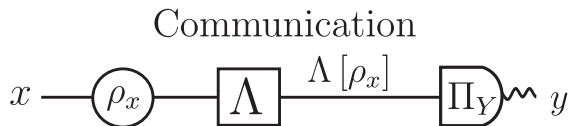
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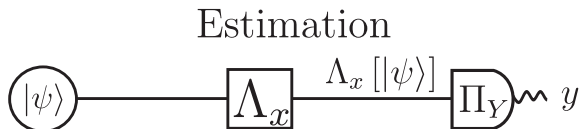
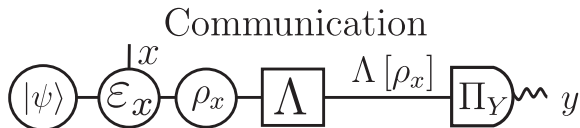
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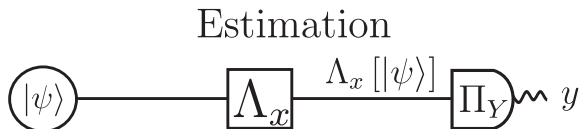
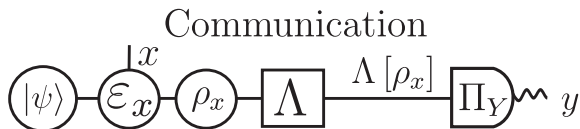
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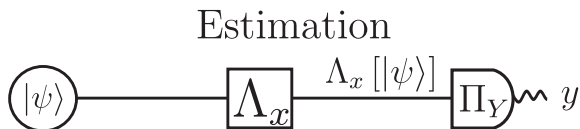
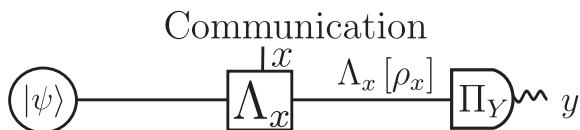
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- Expansion of relative entropy $D(p(y|x) || p(y|x + \delta x)) \approx \frac{F(x)}{2 \ln 2} \delta x^2$,
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- $I(X, Y) = \int p(x) D(p(y|x)||p(y)) dy$.

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$$\chi \approx \frac{\sigma^2 J(\bar{x})}{2 \ln 2} - \sum_{n=r+1} \frac{\sigma^2 F_n(\bar{x})}{4} \log_2 \frac{\sigma^2 F_n(\bar{x})}{4e}$$

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- Obvious: maximize Fisher information $C^{(1,1)} \approx \frac{\sigma^2}{2 \ln 2} F_Q(\bar{x})$.
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$$\rho_x \approx \rho_{\bar{x}} + (x - \bar{x}) \dot{\rho}_{\bar{x}} + \frac{(x - \bar{x})^2}{2} \ddot{\rho}_{\bar{x}}$$
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- **Warning:** Do not increase k to much! $\sigma^2 \ll 1/F(\bar{x})$

Thermal Channel

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- Send coherent states $|\alpha\rangle = D(\alpha)|0\rangle \rightarrow$ encoding in the x quadrature.

Thermal Channel

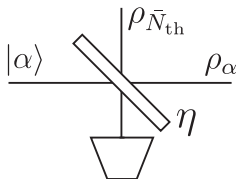
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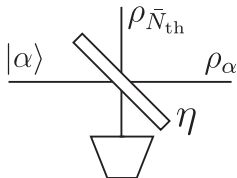
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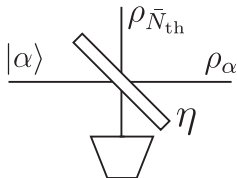


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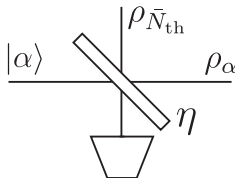
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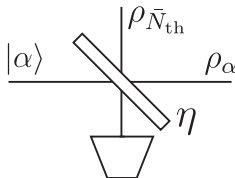
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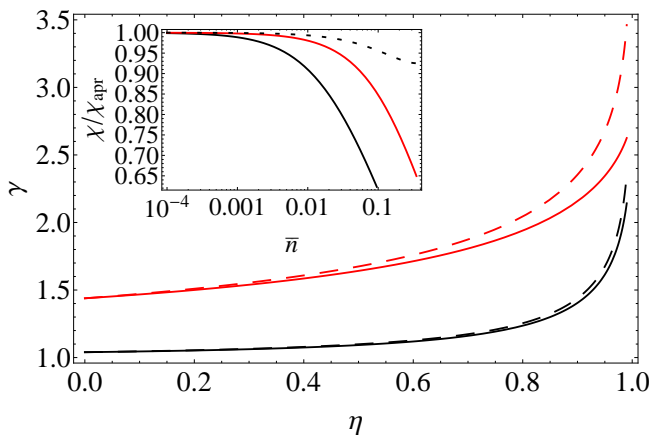


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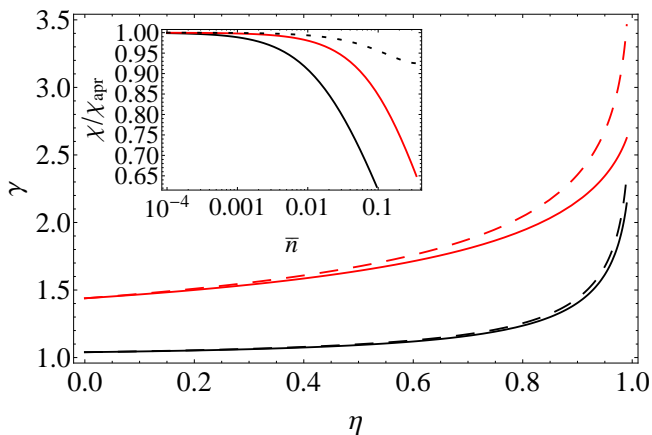
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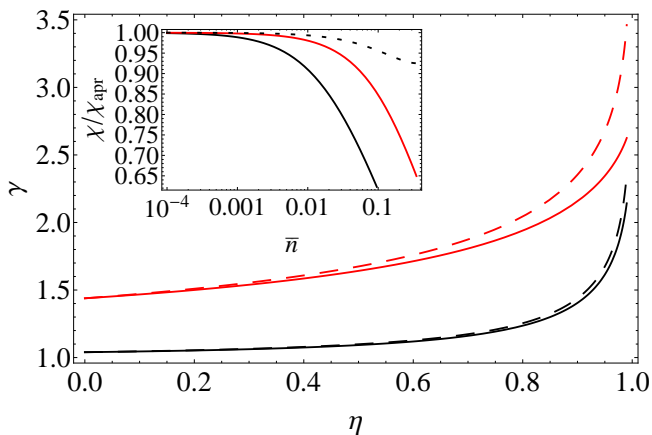


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- Inset - convergence of the approximation: $\bar{N}_{\text{th}} = 0$, $\eta = 0.9$ - dotted; $\bar{N}_{\text{th}} = 0.1$, $\eta = 0.5$ - red; $\bar{N}_{\text{th}} = 1$, $\eta = 0.99$ - black, solid.

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